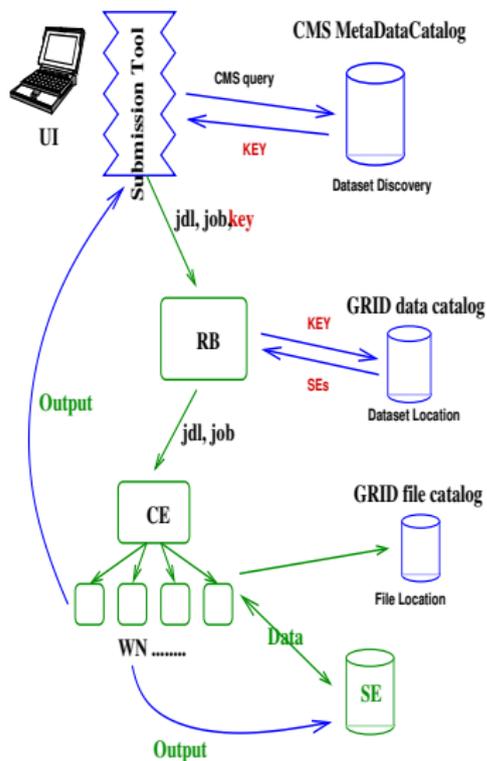
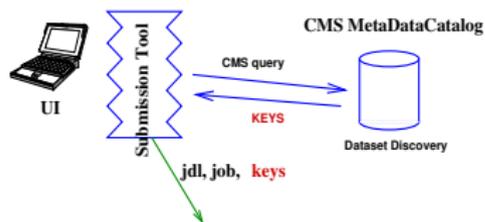


Catalogs: proposed architecture



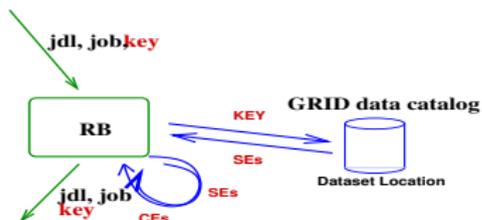
► Three level of catalogs, with defined responsibility and scope

1. **CMS specific Dataset Bookkeeping System**
user access point to next step: CMS responsibility, Centralized (replicated)
2. **Data Location Service**
should be Grid responsibility, **not** CMS, Global (replicated/distributed)
3. **Local File Catalog**
Grid responsibility, Local



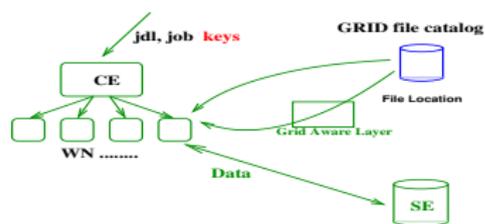
► CMS Dataset Bookkeeping System

- User (or user oriented tool CRAB) access point to Data Discovery and access
- Input user Query: any type, possibly Google like
- Must know about all available data, together with all Data attributes (sw version, calibration, detector condition, processing cards, etc. . .)
- **Does not know about data location**
- Return list of keys corresponding to Data Blocks
- Key list will be passed to next catalog



► Data Location Service

- Grid responsibility
- Accessed at Resource Broker level: Global
- Accessible also from UI
- Input is list of keys, corresponding to Data Blocks
- `inputData = 'key1', 'key2', ...`
- Output is list of Storage Elements hosting Data Blocks
- **Not direct files knowledge is needed**
- Data Discovery is done using DB or collection of Data Blocks
- RB finds CEs fine for SEs and choose CE
- keys are (eventually) sent further down



► Local File Catalog

- Accessed at CE/WN level: Local
- Input is again list of keys, Data Blocks
- Output is list of physical files corresponding to required Data Blocks
- Directly used by COBRA, if POOL file catalog
- Or transformed into POOL format by Grid aware layer (CRAB job wrapper)
- If one POOL catalog (mysql) per site, do not need to extract Data Block fragment: COBRA application uses only what is needed

Feature and requirements

- ▶ **Must answer to question: where is the data**
- ▶ Data is defined in term of “Data Block”, **not as list of files**
- ▶ Data Block Definition is fully experiment specific
- ▶ Data Block does correspond to a set of files, but this information is not needed at DLS level
- ▶ Data Block composition is static
- ▶ DLS get as input a set of keys which correspond to DataBlock, and should return the list of SE which contains this Data Block
- ▶ Interface could be DLI like (or anything else)
- ▶ Possible to query the service directly from UI



- ▶ A Data Block is “atomic”: unbreakable. Don't need to deal with Data Block fragment
- ▶ No or very limited MetaData
- ▶ Definition of Data Block and attribute of Data Block are dealt with in the Dataset Bookkeeping System
- ▶ Eventually MetaData only for Data availability (such as: Data Block is on Tape, on Disk, pinned, pre-staged, etc . . .)
- ▶ Custodianship of the Data (“master” copy, never to be deleted)
- ▶ Data Block can be hierarchical: Dataset is made of Data Blocks, still under discussion
- ▶ Dataset composition can evolve



- ▶ Filling the DB is Experiment responsibility.
- ▶ **When new data is created**, an entry is added into Dataset Bookkeeping System describing all the attribute of the Data and the “key”
- ▶ Uniqueness of Keys is guaranteed by DBS
- ▶ To allow different scope for data (CMS, Analysis group, single person), could foresee a “namespace” for the key, eventually with a naming convention (CMS::key1, CMS:PRSMU:key2, etc. . .)
- ▶ When data is actually created, the initial location of Data Block is registered into the DLS
- ▶ Registration must be Transactional
- ▶ **When data is replicated** a new entry in the DLS is added by the data movement tool (Phedex) when the copy is complete and validated
- ▶ **When data is removed** the corresponding entry in the DLS is removed, then the data is removed.

